

# ALGORITMOS DE ORDENACION EN MEMORIA PRINCIPAL

	Que Hace	Coste mejor caso	Coste peor caso	Coste medio
Inserción directa	comp.	$C_{\min} = n - 1$	$C_{\max} = \sum_{i=1}^n i - 1 = \frac{n^2 + n - 2}{2}$	$C_{prom} = \frac{1}{2} C_{\max} = \frac{n^2 + n - 2}{4}$
	Móv.	$M_{\min} = 2(n - 1)$	$M_{\max} = \frac{n^2 + 9n - 10}{2}$	$M_{prom} = \frac{n^2 + 3n - 4}{4}$
Inserción binaria	comp.	$C_{\min} = n - 1$	$\Omega(I.B.) = \Theta(n^2)$	$C = n(\log_2 n - \log_2 e) + \log_2$
	Móv.	$M_{\min} = 2(n - 1)$	$M_{\max} = \frac{n^2 + 9n - 10}{2}$	$M_{prom} = \frac{n^2 + 3n - 4}{4}$
Selección directa	comp.	$C = 1 + 2 + \dots + n - 1 = \frac{n^2 - n}{2}$		
	Móv.	$3(n - 1)$	$M_{\max(jf)} = 1 + 3 + 5 + \dots + \frac{n}{2}$ $M_{\max} = 3(n - 1) + \frac{n^2}{4}$	$M_{prom} \cong n(\ln n + \gamma)$ $\gamma = 0,577216\dots$
Intercambio directo	comp.	$C = 1 + 2 + \dots + n - 1 = \frac{n^2 - n}{2}$		
	Móv.	$M_{\min} = 0$	$M_{\max} = 3C = 3 \cdot \frac{(n^2 - n)}{2}$	$M_{prom} = \frac{1}{2} M_{\max} = 3 \cdot \frac{(n^2 - n)}{4}$
Sacudida	comp.	$C_{\min} = n - 1$	$C_{\max} = n - k_i \sqrt{n}$	$C_{prom} = \frac{1}{2} (n^2 - n(k_2 + \ln n))$
	Móv.	$M_{\min} = 0$	$M_{\max} = 3C = 3 \cdot \frac{(n^2 - n)}{2}$	$M_{prom} = \frac{1}{2} M_{\max} = 3 \cdot \frac{(n^2 - n)}{4}$
Shell	observaciones	No se explica		
	Coste asintótico siendo n potencia de 2	$O\left(\sum_{i=1}^{\frac{n}{2}} i - 1\right) \equiv \Theta(n^2)$		$O\left(\sum_{i=1}^t \frac{n^2}{h_i}\right) \equiv \Theta(n^2)$
Montón	comp.	$C_{\min} = 2n(\log i)$	$C_{\max} = 2n \log n - O(n)$	$C_{prom} = 2n \log n - O(n \log(\log n))$
	Móv.			
Quicksort	observaciones	Algoritmo de naturaleza recursiva		
	Coste asintótico	$T(n) = 2T\left(\frac{n}{2}\right) + cn$ $= O(n \log n)$	$T(n) = T(n - 1) + cn, \quad n > 1$ $T(n) = T(1) + c \sum_{i=2}^n i = O(n^2)$	$T(n) = O(n \log n)$